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## **THE FORCING TOTAL MONOPHONIC NUMBER OF A GRAPH**

**Abstract:** For a connected graph  $G = (V, E)$  of order at least two. A total monophonic set of a graph  $G$  is a monophonic set  $M$  such that the subgraph induced by  $M$  has no isolated vertices. The minimum cardinality of a total monophonic set of  $G$  is the total monophonic number of  $G$  and is denoted by  $m_t(G)$ . A subset  $T$  of a minimum total monophonic set  $M$  of  $G$  is a forcing total monophonic subset for  $M$  if  $M$  is the unique minimum total monophonic set containing  $T$ . A forcing total monophonic subset for  $M$  of minimum cardinality is a minimum forcing total monophonic subset of  $M$ . The forcing total monophonic number  $fm_t(M)$  in  $G$  is the cardinality of a minimum forcing total monophonic subset of  $M$ . The forcing total monophonic number of  $G$  is  $fm_t(G) = \min\{fm_t(M)\}$ , where the minimum is taken over all minimum total monophonic sets  $M$  in  $G$ . It is shown that for every pair  $a, b$  of positive integers with  $0 \leq a < b$  and  $b > 2a + 1$ , there exists a connected graph  $G$  such that  $fm_t(G) = a$  and  $m_t(G) = b$

**Keywords:** forcing monophonic number, forcing monophonic set, forcing total monophonic number, Monophonic set, monophonic number, total monophonic set, total monophonic number.

**AMS Subject classification:** 05C12.

### **I. INTRODUCTION**

By a graph  $G = (V, E)$ , we mean a finite undirected connected graph without loops or multiple edges. The order and size of  $G$  are denoted by  $p$  and  $q$  respectively. For basic graph theoretic terminology we refer to Harary[2]. The distance  $d(u, v)$  between two vertices  $u$  and  $v$  in a connected graph  $G$  is the length of a shortest  $u$ - $v$  path of length  $d(u, v)$  is called a  $u$ - $v$  geodesic. A vertex  $v$  is said to lie on an  $x$ - $y$  geodesic  $P$  if  $v$  is a vertex of  $P$  including the vertices  $x$  and  $y$ . A set  $S$  of vertices is a geodetic set if  $I[S] = V$ , and the minimum cardinality of a geodetic set is the geodetic number  $g(G)$ . A geodetic set of cardinality  $g(G)$  is called a  $g$ -set. The geodetic number of a

graph was introduced in [3]. The neighborhood of a vertex  $v$  is the set  $N(v)$  consisting of all vertices  $u$  which are adjacent with  $v$ . The closed neighbourhood of a vertex  $v$  is the set  $N(v) = N(v) \cup \{v\}$ . A vertex  $v$  is an extreme vertex if the subgraph induced by its neighbours is complete. A vertex  $v$  is a semi-extreme vertex of  $G$  if  $\Delta(\langle N(v) \rangle) = |N(v)| - 1$ . In particular, every extreme vertex is a semi - extreme vertex and a semi - extreme vertex need not be an extreme vertex.

A chord of a path  $u_1, u_2, \dots, u_k$  in  $G$  is an edge  $u_i u_j$  with  $j \geq i + 2$ . A path  $P$  is called a monophonic path if it is a chordless path. A set  $M$  of vertices is a monophonic set if every vertex of  $G$  lies on a monophonic path joining some pair of vertices in  $M$ , and the minimum cardinality of a monophonic set is the monophonic number  $m(G)$ . The monophonic number of a graph  $G$  was studied in [6]. A set  $S$  of vertices of a graph  $G$  is an edge monophonic set if every edge of  $G$  lies on an  $x$ - $y$  monophonic for some elements  $x$  and  $y$  in  $S$ . The minimum cardinality of an edge monophonic set of  $G$  is the edge monophonic number of  $G$ , denoted by  $em(G)$ . The edge monophonic number of a graph  $G$  was studied in [5].

A total monophonic set of a graph  $G$  is a monophonic set  $M$  such that the subgraph induced by  $M$  has no isolated vertices. The minimum cardinality of a total monophonic set of  $G$  is the total monophonic number of  $G$  and is denoted by  $m_t(G)$ . A total edge monophonic set of a graph  $G$  is an edge monophonic set  $M$  such that the subgraph induced by  $M$  has no isolated vertices. The minimum cardinality of a total edge monophonic set of  $G$  is the total edge monophonic number of  $G$  and is denoted by  $em_t(G)$ . The total edge monophonic number of a graph was introduced and studied in [1].

**Example 1.1.** For the graph  $G$  given in Figure 1.1,  $M_1 = \{v_1, v_2, v_4, v_5\}$ ,  $M_2 = \{v_1, v_2, v_5, v_6\}$ ,  $M_3 = \{v_1, v_2, v_3, v_6\}$ ,  $M_4 = \{v_1, v_3, v_4, v_5\}$  are the minimum total monophonic sets of  $G$  and so  $m_t(G) = 4$ .

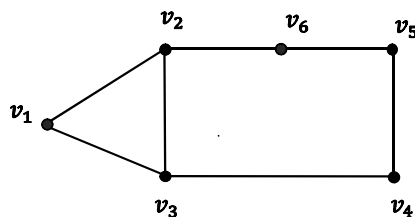


Figure 1.1

Let  $G$  be a connected graph and  $M$  be a minimum monophonic set of  $G$ . A subset  $T \subseteq M$  is called a forcing subset for  $M$  if  $M$  is the unique minimum monophonic set containing  $T$ . A forcing subset for  $M$  of minimum cardinality is a minimum forcing subset of  $M$ . The forcing monophonic number of  $G$ , denoted by  $f_m(G)$ , is  $f_m(G) = \min \{|M| : M \text{ is a minimum monophonic set of } G \text{ and } M \text{ has a forcing subset of size } k\}$  where the minimum is taken over all minimum monophonic sets  $M$  in  $G$ . The forcing monophonic number of a graph was introduced and studied in [4].

A connected graph  $G$  may contain more than one minimum total monophonic sets. For example, the graph  $G$  given in Figure 1.1 contains four minimum total monophonic sets. For each minimum total monophonic set  $M$  in  $G$  there is always some subset  $T$  of  $M$  that uniquely determines  $M$  as the minimum total monophonic set containing  $T$ . Such sets are called forcing total monophonic subsets.

The following theorems used in the sequel.

**Theorem 1.1 [1].** Each extreme vertex and each support vertex of a connected graph  $G$  belongs to every total edge monophonic set of  $G$ . If the set  $M$  of all extreme vertices and support vertices form a total edge monophonic set, then  $M$  is the unique minimum total edge monophonic set of  $G$ .

**Theorem 1.2[1].** For the complete graph  $K_p$ ,  $em_t(K_p) = p$ .

**Theorem 1.3 [1].** For any non trivial tree  $T$ , the set of all end vertices and support vertices of  $T$  is the unique minimum total edge monophonic set of  $G$ .

**Theorem 1.4 [1].** For any connected graph  $G$ ,  $em_t(G) = 2$  if and only if  $G = K_2$ .

**Theorem 1.5 [1].** No cutvertex of a connected graph  $G$  belongs to any minimum total edge monophonic set of  $G$ .

**Theorem 1.6[4].** For the complete graph  $K_p$ ,

Throughout this paper  $G$  denotes a connected graph with at least two vertices.

## 2. FORCING TOTAL MONOPHONIC NUMBER OF A GRAPH

**Definition 2.1.** Let  $G$  be a connected graph and let  $M$  be a minimum total monophonic set of  $G$ . A subset  $T$  of a minimum total monophonic set  $M$  of  $G$  is a forcing total monophonic subset for  $M$  if  $M$  is the unique minimum total monophonic set containing  $T$ . A forcing total monophonic subset for  $M$  of minimum cardinality is a minimum forcing total monophonic subset of  $M$ . The forcing total monophonic number  $fm_t(M)$

in  $G$  is the cardinality of a minimum forcing total monophonic subset of  $M$ . The forcing total monophonic number of  $G$  is  $fm_t(G) = \min\{fm_t(M)\}$ , where the minimum is taken over all minimum total monophonic sets in  $G$ .

**Example 2.2.** For the graph  $G$  given in Figure 2.1,  $M_1 = \{u_1, u_3, u_6, u_7\}$ ,  $M_2 = \{u_1, u_2, u_6, u_7\}$ ,  $M_3 = \{u_1, u_2, u_5, u_6\}$ ,  $M_4 = \{u_1, u_3, u_5, u_6\}$  are the minimum total monophonic sets of  $G$ . It is clear that  $fm_t(M_1) = 2$ ,  $fm_t(M_2) = 2$ ,  $fm_t(M_3) = 2$ ,  $fm_t(M_4) = 2$ , so that  $fm_t(G) = 2$ .

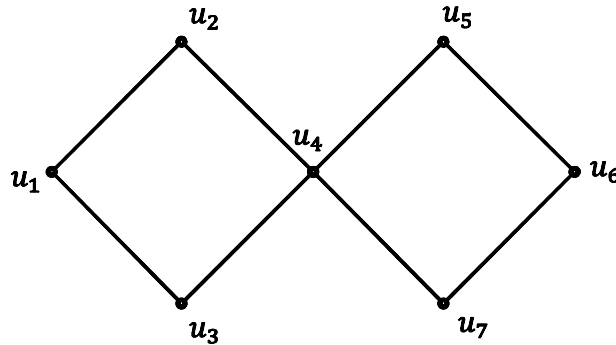


Figure 2.1

**Theorem 2.3.** For a connected graph  $G$ ,  $0 \leq fm_t(G) \leq m_t(G) p$ .

**Proof.** It is clear from the definition of  $fm_t(G)$  that  $fm_t(G) \geq 0$ . Let  $M$  be a minimum total monophonic set of  $G$ . Since  $fm_t(M) \leq m_t(G)$  and since  $fm_t(G) = \min\{fm_t(M) : M \text{ is a minimum total monophonic set in } G\}$ , it follows that  $fm_t(G) \leq m_t(G)$ . Thus  $0 \leq fm_t(G) \leq m_t(G) \leq p$ .

**Remark 2.4.** The bounds in Theorem 2.3 are sharp. By Theorems 1.2 and 1.6, for the complete graph  $Kp$  ( $p \geq 2$ ),  $m_t(Kp) = p$ , also  $V(Kp)$  is the unique total monophonic set of  $Kp$  and so  $fm_t(Kp) = 0$ . The inequalities in Theorem 2.3 are strict.

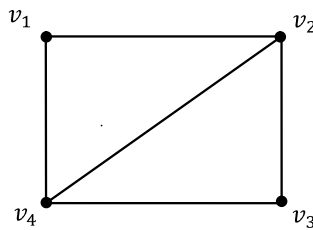


Figure 2.2

For the graph  $G$  given in Figure 2.2,  $M_1 = \{v_1, v_2, v_3\}$ ,  $M_2 = \{v_1, v_3, v_4\}$  are the minimum total monophonic sets of  $G$  so that  $m_t(G) = 3$ . It is clear that  $fm_t(M_1) = 1$  and  $fm_t(M_2) = 1$  so that  $fm_t(G) = 1$ . Thus  $0 < fm_t(G) < m_t(G) < p$ .

**Theorem 2.5.** Let  $G$  be a connected graph. Then

- (i)  $fm_t(G) = 0$  if and only if  $G$  has a unique minimum total monophonic set.
- (ii)  $fm_t(G) = 1$  if and only if  $G$  has at least two minimum total monophonic sets, one of which is a unique minimum total monophonic set containing one of its elements, and
- (iii)  $fm_t(G) = m_t(G)$  if and only if no minimum total monophonic set of  $G$  is the unique minimum total monophonic set containing any of its proper subsets.

**Proof.**(i) Let  $fm_t(G) = 0$ . Then by definition,  $fm_t(M) = 0$  for some minimum total monophonic set  $M$  of  $G$ , so that the empty set  $\emptyset$  is the minimum forcing subset for  $M$ . Since the empty set  $\emptyset$  is a subset of every set, it follows that  $M$  is the unique minimum total monophonic set of  $G$ . The converse is clear.

(ii) Let  $fm_t(G) = 1$ . Then by Theorem 2.5 (i),  $G$  has atleast two minimum total monophonic sets. Since  $fm_t(G) = 1$ , there is a singleton subset  $T$  of a minimum total monophonic set  $M$  of  $G$  such that  $T$  is not a subset of any other minimum total monophonic set of  $G$ . Thus  $M$  is the unique minimum total monophonic set containing one of its elements. The converse is clear.

(iii) Let  $fm_t(G) = m_t(G)$ . Then  $fm_t(M) = m_t(G)$  for every minimum total monophonic set in  $G$ . By Theorem 2.3,  $m_t(G) \geq 2$ ,  $fm_t(G) \geq 2$ . Then by Theorem 2.5(i),  $G$  has at least two minimum total monophonic set of  $G$ . Since  $fm_t(M) = m_t(G)$ , no proper subset of  $M$  is a forcing subset of  $M$ . Thus no minimum total monophonic set of  $G$  is the unique minimum total monophonic set containing any of its proper subsets. Conversely,  $G$  contains more than one minimum total monophonic set and no subset of any minimum total monophonic set  $M$  other than  $M$  is a subset for  $M$ . Hence it follows that  $fm_t(G) = m_t(G)$ .

**Definition 2.6.** A vertex  $v$  of a connected graph  $G$  is said to be a total monophonic vertex of  $G$  if  $v$  belongs to every minimum total monophonic set of  $G$ .

**Example 2.7.** For the graph  $G$  given in Figure 2.3,  $M_1 = \{v_1, v_4, v_5, v_8\}$ ,  $M_2 = \{v_1, v_2, v_4, v_5\}$  are the only two total monophonic sets of  $G$ . It is clear that  $v_1, v_5$  are total monophonic vertices of  $G$ .

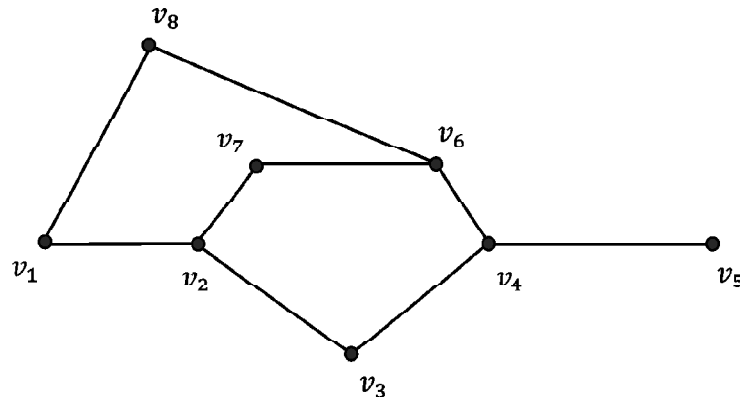


Figure 2.3

**Corollary 2.8.** Let  $G$  be a connected graph and let  $M$  be a minimum total monophonic set of  $G$ . Then no total monophonic vertex of  $G$  belongs to any minimum forcing total monophonic subset of  $M$ .

**Theorem 2.9.** Let  $G$  be a connected graph and let  $S$  be the set of all total monophonic vertices of  $G$ . Then  $fm_t(G) \leq m_t(G) - |S|$ .

**Proof.** Let  $M$  be any minimum total monophonic set of  $G$ . Then  $m_t(G) = |M|$ ,  $S \subseteq M$  and  $M$  is the unique minimum total monophonic set containing  $M - S$ .

Thus  $fm_t(G) \leq |M - S| = |M| - |S| = m_t(G) - |S|$ .

**Corollary 2.10.** If  $G$  is a connected graph with  $l$  extreme vertices and  $k$  support vertices, then  $fm_t(G) \leq m_t(G) - (l+k)$ . **Proof.** This follows from Theorem 1.1 and Theorem 2.9.

**Remark 2.11.** The bound in Theorem 2.9 is sharp. For the graph  $G$  given in Figure 2.2,  $m_t(G) = 3$  and  $fm_t(G) = 1$ .

Also  $S = \{v_1, v_3\}$  is the set of all total monophonic vertices of  $G$  and  $fm_t(G) = m_t(G) - |S|$ . Also the inequality in Theorem 2.9 can be strict.

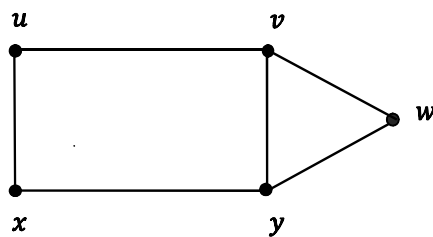


Figure 2.4

For the graph  $G$  given in Figure 2.4,  $M_1 = \{u, v, w\}$ ,  $M_2 = \{x, y, w\}$  are the minimum total monophonic sets of  $G$  and so that  $m_t(G) = 3$ . Since  $M_1$  is the unique minimum total monophonic set contains the subset  $\{u\}$  so that  $fm_t(M_1) = 1$  and  $M_2$  is the unique minimum total monophonic set contains the subset  $\{x\}$  so that  $fm_t(M_2) = 1$ . Hence, we have  $fm_t(G) = 1$ . Also, the vertex  $w$  is the unique total monophonic vertex of  $G$ , we have  $fm_t(G) < m_t(G) |S|$ .

**Theorem 2.12.** Let  $G$  be a connected graph and let  $M$  be a minimum total monophonic set of  $G$ . Then no cut vertex of  $G$  belongs to any minimum forcing total monophonic subset of  $M$ .

**Proof:** Let  $v$  be a cut vertex of  $G$  which is not a support vertex. By Theorems 1.1 and 1.5,  $v$  does not belong to any minimum total monophonic set of  $G$ . Since any minimum forcing total monophonic subset of  $M$  is a subset of minimum total monophonic set. It follows that, no cut vertex of  $G$  belongs to any minimum forcing total monophonic subset of  $M$ .

**Theorem 2.13:** For any complete graph  $G = K_p (p \geq 2)$  or any non trivial tree  $G = T$ ,  $fm_t(G) = 0$ .

**Proof:** For  $G = K_p$ , it follows from Theorem 1.6 that the set of all vertices of  $G$  is the unique minimum total monophonic set of  $G$ . Now, it follows from Theorem 2.5 (i) that  $fm_t(G) = 0$ . If  $G$  is a non trivial tree, then by Theorem 1.3, the set of all end vertices and support vertices of  $G$  is the unique minimum total monophonic set of  $G$  and by Theorem 2.5(i)  $fm_t(G) = 0$ .

**Theorem 2.14:** If  $G$  is a connected graph with  $m_t(G) = 2$ , then  $fm_t(G) = 0$ .

**Proof:** This follows from Theorem 1.4 and Theorem 2.13.

**Theorem 2.15:** For the complete bipartite graph  $G = (m, n \geq 2)$ ,  $fm_t(G) =$

$$\left\{ \begin{array}{l} 1 \text{ if } 2 = m = n \\ 4 \text{ if } 3 \leq m \leq n \\ 1 \text{ if } 2 = m < n \\ 0 \text{ if } m = 1 \text{ and } n = 2 \end{array} \right\}$$

**Proof:** We prove this theorem by considering four cases.

Case 1: If  $m=1$  and  $n=2$ . Let  $U = \{u_1, u_2, \dots, u_m\}$  and  $W = \{w_1, w_2, \dots, w_n\}$  be the bipartition of  $G$ . If  $m=1$ , then  $G$  is a tree and its forcing total monophonic number is 0.

Case 2: If  $2=m=n$ , then  $U = \{u_1, u_2\}$  and  $W = \{w_1, w_2\}$  and any minimum total monophonic set of  $G$  is of the following forms: (i)  $U \cup \{w_i\}$  or (ii)  $W \cup \{u_i\}$ . Hence in both cases  $fm_t(G) = 1$ .

Case 3: If  $2=m < n$ , then for any  $j(1 \leq j \leq n)$ ,  $S_j = U \cup \{w_j\}$  is a minimum total monophonic set of  $G$ . Since  $n \geq 3$ , then by Theorem 2.5 (ii), we have  $fm_t(G) = 1$ .

Case 4: If  $3=m=n$ , then any minimum total monophonic set of  $G$  is of the following forms: (i)  $U \cup \{w_j\}$  for some  $j(1 \leq j \leq n)$ . (ii)  $W \cup \{u_i\}$  for some  $i(1 \leq i \leq m)$ , or (iii) any set got by choosing any two elements from each of  $U$  and  $W$ . If  $3=m < n$ , then any minimum total monophonic set of  $G$  is either  $U \cup \{w_j\}$  for some  $j(1 \leq j \leq n)$  or any set got by choosing any two elements from each of  $U$  and  $W$ . Hence in both cases, we have  $m_t(G) = 4$ . Clearly, no minimum total monophonic set of  $G$  is the unique minimum total monophonic set containing any of its proper subsets. Then by Theorem 2.5 (iii), we have  $fm_t(G) = m_t(G) = 4$ .

### 3. REALIZATION RESULTS

**Theorem 3.1:** For every pair  $a, b$  of positive integers with  $a \leq b$ , there exists a connected graph  $G$  such that  $fm_t(G) = a$  and  $m_t(G) = b$ .

**Proof:** If  $a=0$ , let  $G = K_b$ . Then by Theorem 2.13,  $fm_t(G) = 0$  and by Theorem 1.2,  $m_t(G) = b$ . Thus we assume that  $0 < a < b$ .

For each  $i$  with  $1 \leq i \leq a$ , let  $C_i : u_{i,1}, u_{i,2}, u_{i,3}, u_{i,4}, u_{i,1}$  be a cycle of order 4. Let  $K_{1,b-2a-1}$  be a star with the cutvertex  $x$  and  $V(K_{1,b-2a-1}) = \{x, v_1, v_2, \dots, v_{b-2a-1}\}$ . The graph  $G$  is obtained from  $C_i(1 \leq i \leq a)$  and  $K_{1,b-2a-1}$  by joining the vertices  $x$  and  $u_{i,1}$  and by joining the vertices  $u_{i,2}$  and  $u_{i,4}(1 \leq i \leq a)$ . The graph  $G$  is given in Figure 3.1.

Let  $S = \{v_1, v_2, \dots, v_{b-2a-1}, u_{1,3}, u_{2,3}, \dots, u_{a,3}, x\}$  be the set of all extreme vertices and support vertex of  $G$ . By Theorem 1.1, every total monophonic set of  $G$  contains  $S$ . It is easily verified that  $S$  is not a total monophonic set of  $G$ . Hence  $M' = S \cup \{u_{1,2}, u_{2,2}, \dots, u_{a,2}\}$  is a total monophonic set of  $G$ . Thus  $m_t(G) = b$ .



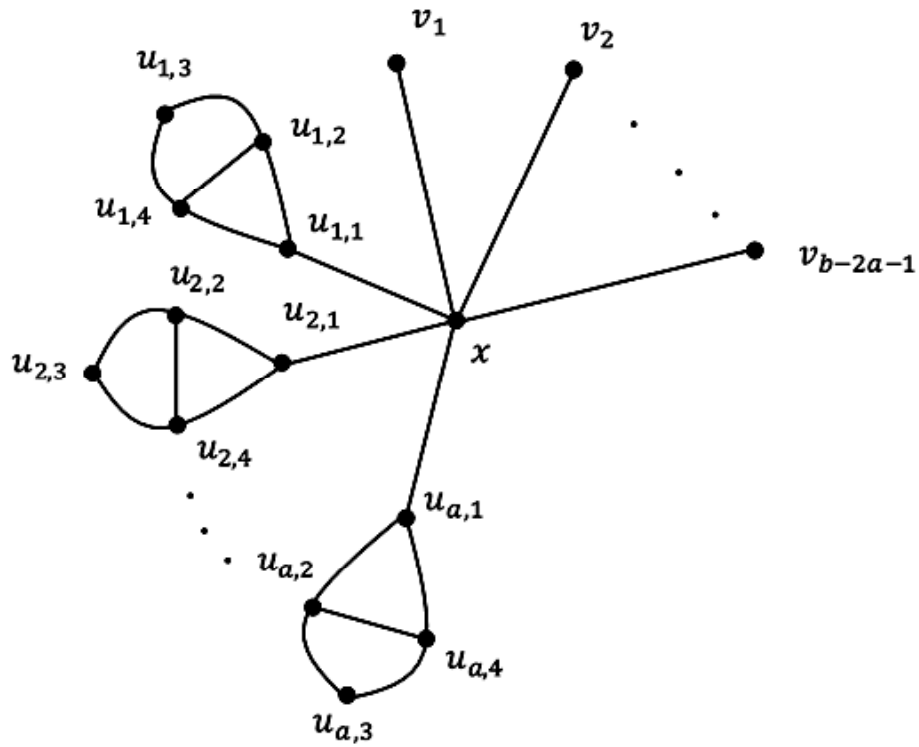


Figure 3.1

Next we show that  $fm_t(G) = a$ . It is observed that  $S$  is the set of all total monophonic vertices of  $G$ . Then by Theorem 2.9,  $fm_t(G) \leq m_t(G) - a = b - (b - a) = a$ .

Now, since  $m_t(G) = b$  and every total monophonic set of  $G$  contains  $S$ , it is easily seen that every minimum total monophonic set  $M_1$  of  $G$  is of the form  $S \cup \{x_1, x_2, \dots, x_a\}$ , where  $x_i \in \{u_{i,2}, u_{i,4}\}$  for every  $i(1 \leq i \leq a)$ . Let  $T$  be any proper subset of  $M_1$  with  $|T| < a$ . Then there is a vertex  $x \in M_1 - S$  such that  $x \notin T$ . If  $x = u_{i,2}$  ( $1 \leq i \leq a$ ), then  $M_2 = (M_1 - \{u_{i,2}\}) \cup \{u_{i,4}\}$  is a minimum total monophonic set of  $G$  containing  $T$ .

Similarly, if  $x = u_{i,4}$  ( $1 \leq i \leq a$ ), then  $M_3 = (M_1 - \{u_{i,4}\}) \cup \{u_{i,2}\}$  is a minimum total monophonic set of  $G$  containing  $T$ . Thus  $M_1$  is not the unique minimum total monophonic set of  $G$  containing  $T$  and so  $T$  is not a forcing total monophonic subset of  $M_1$ . Since this is true for all minimum total monophonic set of  $G$ , it follows that  $fm_t(G) \geq a$  and so  $fm_t(G) = a$ .

**Theorem 3.2:** For every integers  $a$  and  $b$  with  $a < b$ , and  $b - 2a - 2 > 0$ , there exists a connected graph  $G$  such that,  $fm_t(G) = a$  and  $m_t(G) = b$ .

**Proof:** Case 1.  $a = 0, b \geq 2$ . Let  $G = K_{1,b-1}$ . Then by Theorem 2.13,  $fm_t(G) = 0$  and  $m_t(G) = b$ .

Case 2.  $0 < a < b$ . Let  $F_i: s_i, r_i, u_i, t_i, s_i$  be a copy of  $C_4$ . Let  $H$  be a graph obtained from  $F_i$ 's by identifying  $t_{i-1}$  of  $F_{i-1}$  and  $s_i$  of  $F_i$  ( $2 \leq i \leq a$ ). Let  $G$  be a graph obtained from  $H$  by adding  $b - 2a - 1$  new vertices  $x, v_1, v_2, \dots, v_{b-2a-2}$  and joining the edges  $xs_1, t_a v_1, \dots, t_a v_{b-2a-2}$  as shown in Figure 3.2.

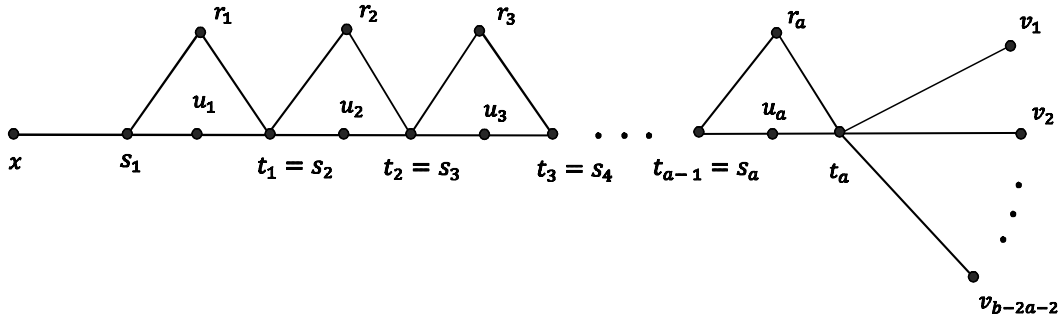


Figure 3.2

Let  $Z = \{x, v_1, v_2, \dots, v_{b-2a-2}\}$  be the set of end vertices of  $G$ . It is clear that  $Z$  is not a total monophonic set of  $G$ . By Theorem 1.5,  $Z' = Z \cup \{s_1, s_2, \dots, s_a, t_a\}$  is a subset of every total monophonic set of  $G$ . Let  $H_i = \{u_i, r_i\}$  ( $1 \leq i \leq a$ ). We observe that every  $m_t(G)$  must contain at least one vertex from each  $H_i$ , so that  $m_t(G) \geq b - 2a - 1 + a + 1 + a = b$ . Now,  $M = Z' \cup \{r_1, r_2, \dots, r_a\}$  is a total monophonic set of  $G$ , so that  $m_t(G) \leq b - 2a - 1 + a + 1 + a = b$ . Thus  $m_t(G) = b$ . Next, we show that  $fm_t(G) = a$ . Since every  $m_t$ -set contains  $Z'$ , it follows from Theorem 2.9,  $fm_t(G) \leq m_t(G) - (b - 2a - 1 + a + 1) = a$ .

It is easily seen that every  $m_t$ -set of  $G$  is of the form  $Z' \cup \{r_1, r_2, \dots, r_a\}$ , where  $r_i \in H_i$  ( $1 \leq i \leq a$ ). Let  $T$  be any proper subset of  $M$  with  $|T| < a$ . Then there exists  $r_i$  ( $1 \leq i \leq a$ ) such that  $r_i \notin T$ . Let  $e_i$  be the vertex of  $H_i$  distinct from  $r_i$ . Then  $W = (M - \{r_i\}) \cup \{e_i\}$  is  $am_t$ -set properly containing  $T$ . Thus  $M$  is not the unique  $m_t$ -set containing  $T$ , so that  $T$  is not a forcing subset of  $M$ . This is true for all  $m_t$ -sets, so that  $fm_t(G) = a$ .

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