

## C-ALGEBRA AS A POSET

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**ABSTARCT:** *In this paper we define a partial ordering  $\leq$  on a C-Algebra and prove that the centre  $B(A)$  of a C-algebra  $A$  with  $T$  becomes a Boolean algebra under this partial ordering. We derive a number of necessary and sufficient conditions for a C-algebra  $\langle A, \vee, \wedge, \prime \rangle$  to become a Boolean algebra.*

### INTRODUCTION

Kleene [3] introduced the notion of a regular extension of classical logic. For our purpose a regular extension of an algebra  $A$  is obtained by adjoining a new element  $U$  (for “un Known”, “undetermined” or undefined”) to  $A$  and extending the operations on  $A$  to  $A \cup \{U\}$ . If  $A$  is the two element Boolean algebra  $B = \{T, F\}$ , then  $A \cup \{U\} := \{T, F, U\}$  is denoted by  $C$ .

Guzman and squier [2] introduced the variety of C-algebras as the variety generated by this 3-element algebra  $C = \{T, F, U\}$  which is the algebraic form of the three valued conditional logic. U.M. Swamy, G.C. Rao and R.V.G Ravi Kumar [6] introduced the concept of the centre of a C-algebra  $A$  which is denoted by  $B(A)$  and proved that the centre of a C-algebra is a Boolean algebra. Many mathematicians like Berman. J, McCarthy. J, Kleene. S have worked on three valued logic.

In this paper we introduce a partial ordering  $\leq$  on a C-Algebra and study the properties of C-algebra as a poset. We prove that the centre  $B(A)$  of a C-algebra  $A$  with  $T$  becomes a Boolean algebra under this partial ordering. We derive a number of necessary and sufficient conditions for a C-algebra  $\langle A, \vee, \wedge, \prime \rangle$  to become a Boolean algebra.

### 1. Preliminaries

In this section we recall the definition of C-algebra and some results from [2] and [6]. Also we prove some more results in a C-algebra which will be required later.

**Definition 1.1 [2]:** By a C-algebra we mean an algebra of type (1,2,2) with operations  $\prime$ ,  $\wedge$  and  $\vee$  satisfying the following identities.

- (a)  $x'' = x$
- (b)  $(x \wedge y) \prime = x \prime \vee y \prime$
- (c)  $(x \wedge y) \wedge z = x \wedge (y \wedge z)$
- (d)  $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$
- (e)  $(x \vee y) \wedge z = (x \wedge z) \vee (x \prime \wedge y \wedge z)$
- (f)  $x \vee (x \wedge y) = x$
- (g)  $(x \wedge y) \vee (y \wedge x) = (y \wedge x) \vee (x \wedge y)$

**Example 1.2 [2]:** The two element algebra  $B = \{T, F\}$  is a C-algebra with operations  $\prime$ ,  $\wedge$  and  $\vee$  defined as in the following table.

X	X'	$\wedge$	T	F	$\vee$	T	F
T	F	T	T	F	T	T	T
F	T	F	F	F	F	T	F

**Example : 1.3 [2] :** The 3 - element algebra  $C = \{T, F, U\}$  is a c - algebra with operations  $\vee$ ,  $\wedge$ ,  $\prime$ , defined as in the following tables.

X	X'	$\wedge$	T	F	U	$\vee$	T	F	U
T	F	T	T	F	U	T	T	T	T
F	T	F	F	F	F	F	T	F	U
U	U	U	U	U	U	U	U	U	U

**Note 1.4 [2]:**

- (i) Identities 1.1a, 1.1b imply that the variety of C-algebras satisfies all the dual statements of 1.1b to 1.1g
- (ii)  $\wedge$  and  $\vee$  are not commutative in C
- (iii) The ordinary right distributive law of  $\wedge$  over  $\vee$  fail in C
- (iv) Every Boolean algebra is a C-algebra

(v) Throughout this paper  $A$  represents a C-algebra unless otherwise mentioned.

**Lemma: 1.5 [2]:** Every C - algebra satisfies the following laws

- (a)  $x \wedge x = x$
- (b)  $x \wedge y = x \wedge (x' \vee y)$
- (c)  $x \vee (x' \wedge x) = x$
- (d)  $(x \vee x') \wedge y = (x \wedge y) \vee (x' \wedge y)$
- (e)  $(x \vee x') \wedge x = x$

**Lemma 1.6 [6]:** Every C - algebra satisfies the laws

- (a)  $x \vee x' = x' \vee x$
- (b)  $x \vee y \vee x = x \vee y$
- (c)  $x \wedge x' \wedge y = x \wedge x'$

**Lemma 1.7 [6]:** If  $a \in A$  then

- (i)  $a$  is left Identity for  $\wedge$  if and only if  $a$  is right Identity for  $\wedge$
- (ii)  $a$  is left identity for  $\vee$  if and only if  $a$  is right identity for  $\vee$

**Definition 1.8 [2]:**  $A$  is said to have T (and F). if there is an Identity for  $\wedge$ . In this case we write T for the Identity for  $\wedge$  and F in place of T'

**Lemma 1.9 [2]:** Let  $A$  be with T and  $x, y \in A$ .

- (a)  $x \vee y = F$  if and only if  $x = y = F$
- (b) if  $x \vee y = T$  then  $x \vee x' = T$

**Lemma: 1.10 [2]:** Every C-algebra with T satisfies the law  $x \vee T = x \vee x'$

**Theorem 1.11[6]:** The following are equivalent in  $A$ .

- (a)  $A$  is a Boolean algebra
- (b)  $x \vee (y \wedge x) = x$  for all  $x, y \in A$
- (c)  $x \wedge y = y \wedge x$  for all  $x, y \in A$
- (d)  $(x \vee y) \wedge y = y$  for all  $x, y \in A$

- (e)  $x \vee x'$  is Identity for  $\wedge$  for all  $x \in A$
- (f)  $x \vee x' = y \vee y'$  for all  $x, y \in A$
- (g)  $A$  has a right zero for  $\wedge$
- (h) For any  $x, y \in A$ , there exists  $a \in A$  such that  $x \wedge a = y \wedge a = a$ .
- (i) For any  $x, y \in A$  if  $x \vee y = y$  then  $y \wedge x = x$

**Definition 1.12 [6]:** Let  $A$  be with  $T$ . An element  $x \in A$  is called a central element of  $A$  if  $x \vee x' = T$ . The set  $\{x \in A \mid x \vee x' = T\}$  of all central elements of  $A$  is called the centre of  $A$  and is denoted by  $B(A)$ .

**Theorem: 1.13:** let  $A$  be with  $T$ . Then  $B(A)$  is a Boolean algebra with induced operations  $\wedge$ ,  $\vee$  and  $'$

Note: 1.14: (i).  $a$  is the identity for  $\wedge$  if and only if  $a'$  is Identity for  $\vee$   
since  $a \wedge x = x = x \wedge a$  for all  $x \in A \Leftrightarrow a' \vee x = x = x \vee a'$  for all  $x \in A$ .

(ii) observe that if Identity for  $\wedge$  exists then it is unique and it is denoted by  $T$ , then  $T'$ , denoted by  $F$ , is the unique Identity for  $\vee$

**Theorem: 1.15:** Let  $a \in A$ . Then the following are equivalent.

- (1)  $a$  is left Identity for  $\wedge$
- (2)  $a$  is right Identity for  $\wedge$
- (3)  $a$  is Identity for  $\wedge$
- (4)  $a'$  is left Identity for  $\vee$
- (5)  $a'$  is right Identity for  $\vee$
- (6)  $a'$  is Identity for  $\vee$

## 2. Some Properties of C-Algebra

Guzman & Squier [2] proved that if  $a \in A$  and  $a' = a$  then  $a$  is left zero for both  $\vee$  and  $\wedge$ . Now we prove the converse.

**Theorem 2.1:** Let  $a \in A$ . Then  $a' = a$  if and only if  $a$  is left zero for both  $\vee$  and  $\wedge$

**Proof:** Suppose that  $a$  is left zero for both  $\vee$  and  $\wedge$

$$a \wedge x = a = a \vee x \text{ for all } x \in A \quad \dots(1)$$

$$\begin{aligned} a' &= a' \vee (a \wedge a') \quad (\text{by Lemma 1.5c}) \\ &= a' \vee a \quad (\text{by (1)}) \\ &= a \vee a' \quad (\text{by Lemma 1.6 (a)}) \\ &= a \quad (\text{by (1)}) \end{aligned}$$

Hence  $a' = a$

In a C-algebra A with T, if  $q = p'$  then  $p \vee q = T$  and  $p \wedge q = F$  need not hold, for example in C,  $U \vee U' = U \wedge U' = U$ . But we prove the converse in the following.

**Theorem 2.2:** Let A be a C - algebra with T and  $p, q \in A$ .

$$\text{If } p \wedge q = F, p \vee q = T \text{ then } q = p'$$

**Proof.** Let  $p \wedge q = F, p \vee q = T$

$$\text{Now } p \vee q = T \Rightarrow p \vee p' = T \quad (\text{By 1.9}) \quad \dots(1)$$

$$\begin{aligned} \text{Now } q &= T \wedge q = (p \vee p') \wedge q \quad (\text{By (1)}) \\ &= (p \wedge q) \vee (p' \wedge q) \quad (1.5d) \\ &= F \vee (p' \wedge q) \\ &= (p' \wedge p) \vee (p' \wedge q) \quad (\text{By (1)}) \\ &= p' \wedge (p \vee q) = p' \wedge T = p' \\ &\text{Hence } q = p' \end{aligned}$$

By 1.5 b, we know that in a C-algebra,  $x \wedge y = x \wedge (x' \vee y)$ . Now we prove the following.

**Theorem 2.3** Every C-algebra satisfies the laws.

$$(i) \ x \wedge y = (x' \vee y) \wedge x, \quad (ii) \ x \vee y = (x' \wedge y) \vee x$$

$$\begin{aligned} \text{Proof:} \quad x \wedge y &= x \wedge (x' \vee y) \\ &= (x \wedge x') \vee (x \wedge y) \\ &= (x' \wedge x) \vee (x \wedge y \wedge x) \quad (\text{by lemma 1.6 a,b}) \\ &= (x' \vee y) \wedge x \quad (\text{by def 1.1}) \end{aligned}$$

$$\text{similarly } x \vee y = (x' \wedge y) \vee x$$

**Corollary 2.4:** For any  $x, y \in A$ , the elements  $x, x' \vee y$  commute.

$$\text{That is, (i) } x \wedge (x' \vee y) = (x' \vee y) \wedge x$$

$$\text{(ii) } x \vee (x' \wedge y) = (x' \wedge y) \vee x$$

**Corollary 2.5:** Every C-algebra satisfies.

$$x \wedge (x' \vee x) = (x' \vee x) \wedge x = x \vee (x' \wedge x) = (x' \wedge x) \vee x = x$$

**Note 2.6:** If  $A$  has an element  $a$  with the property  $a' = a$  then it is unique and is denoted by  $U$ .

**Theorem: 2.7.**  $B$  is the only C-algebra of order two.

**Proof:** Let  $A$  be a C-algebra of order two. Let  $A = \{x, y\}$

If  $x' = x$  then  $y' = y$  (by definition of C-algebra) which is not true (See 2.6)

Therefore  $x' = y, A = \{x, x'\}$ . Since  $x \vee x' = x' \vee x$  (by 1.6 (a))

We get  $x \vee y = y \vee x$  for all  $x, y \in A$  Thus  $A$  is a Boolean algebra of order 2

Hence  $B$  is the only C-algebra of order two

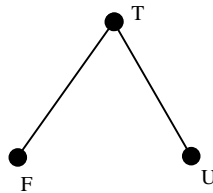
**Theorem: 2.8.** every C-algebra satisfies

$$(x \vee y) \wedge x = x \vee (y \wedge x)$$

### 3. C-ALGEBRA AS A POSET

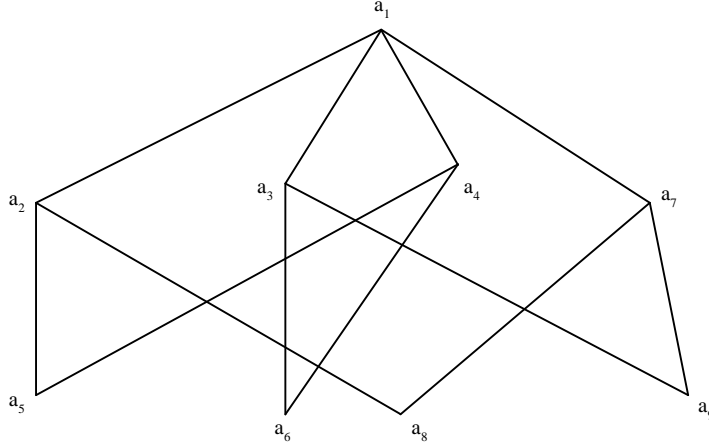
In this section we define a partial ordering “ $\leq$ ” on C-algebra  $A$  and study the properties of the C-algebra as a poset. We prove that the centre  $B(A)$  of a C-algebra  $A$  becomes a Boolean algebra under this partial ordering. We derive a number of necessary and sufficient conditions for a C-algebra  $\langle A, \vee, \wedge \rangle$  to become a Boolean algebra.

**Definition: 3.1** Let  $A$  be a C-algebra. Define ‘ $\leq$ ’ on  $A$  by  $x \leq y$  if  $y \wedge x = x$  then ‘ $\leq$ ’ is a partial ordering on  $C$ .



It is enough to verify that ' $\leq$ ' is antisymmetric, remaining proof is a routine verification. Suppose that  $x, y \in A, x \leq y$  and  $y \leq x$

$$\left\{ \begin{array}{l} y \wedge x = x, x \wedge y = y \end{array} \right\}$$



Now  $x \vee y = x \vee (x \wedge y) = x$   
 $y \vee x = y \vee (y \wedge x) = y$   
 So that  $x = y \wedge x = (y \vee x) \wedge (x \vee y) = (x \vee y) \wedge (y \vee x) = x \wedge y = y$

Hence ' $\leq$ ' is antisymmetric

**Note 3.2:** If a C-algebra A has T. Then T is the greatest element of the poset  $(A, \leq)$  since  $T \wedge x = x$  for all  $x \in A$  (that is  $x \leq T$  for all  $x \in A$ ).

**Example 3.3:** The Hasse diagram of the poset  $(C, \leq)$  is

**Example 3.4:** We know that  $C \times C$  is a C - algebra under point wise operations. The Hasse diagram of the poset  $(C \times C, \leq)$  is given below

$$C \times C \quad a_1 = (T, T), a_2 = (T, F), a_3 = (T, U), a_4 = (F, T), a_5 = (F, F),$$

$$a_6 = (F, U), a_7 = (U, T), a_8 = (U, F), a_9 = (U, U)$$

**Note 3.5 :** The above two examples suggest that, in general, if A is a C-algebra then  $(A, \leq)$  is a join semi lattice. We do not know yet whether  $(A, \leq)$  is a join semi lattice. But we have the following.

**Theorem: 3.6** In the poset  $(A, \leq)$ , for any  $x \in A$

Supremum of  $\{x, x'\} = x \vee x'$ , infimum of  $\{x, x'\} = x \wedge x'$

**Proof:**  $(x \vee x') \wedge x = x$  and  $(x \vee x') \wedge x' = x'$  by cor. 2.5

Therefore  $x \leq x \vee x'$  and  $x' \leq x \vee x'$

Hence  $x \vee x'$  is an upper bound of  $\{x, x'\}$

Suppose  $t$  is an upper bound of  $\{x, x'\}$ .

We have  $t \wedge x = x$ .  $t \wedge x' = x'$

Now  $t \wedge (x \vee x') = (t \wedge x) \vee (t \wedge x') = x \vee x'$

Therefore  $x \vee x'$  is least upper bound of  $\{x, x'\}$ , That is, supremum of  $\{x, x'\} = x \vee x'$ . Similarly, infimum of  $\{x, x'\} = x \wedge x'$

**Corollary: 3.7** In a C-algebra  $x \vee x' = x' \vee x$  and  $x \wedge x' = x' \wedge x$

In general for a C- algebra  $A$  with  $T$ .  $x \vee y$  need not be the lub of  $\{x, y\}$  in  $(A, \leq)$  for example, in  $C$ ,  $U \vee F = U$  is not the lub of  $U$  and  $F$ , However we have the following results.

**Theorem 3.8:** Let  $A$  be a C - algebra with  $T$  and  $x, y \in A$ . If  $x \in B(A)$  and  $x \wedge y = y \wedge x$  then  $\sup \{x, y\} = x \vee y$

**Proof:** Let  $x \wedge y = y \wedge x$  and  $x \vee x' = T$

$$\begin{aligned} (x \vee y) \wedge x &= (x \wedge x) \vee (x' \wedge y \wedge x) = x \vee (x' \wedge x \wedge y) \\ &= x \vee (x \wedge (x' \wedge y)) = x \end{aligned}$$

Therefore  $x \leq x \vee y$

$$\begin{aligned} (x \vee y) \wedge y &= (x \wedge y) \vee (x' \wedge y \wedge y) = (x \wedge y) \vee (x' \wedge y) = (x \vee x') \wedge y \\ &= T \wedge y = y \end{aligned}$$

Therefore  $y \leq x \vee y$ .

Hence  $x \vee y$  is an upper bound of  $\{x, y\}$

Suppose  $u$  is an upper bound of  $\{x, y\}$ . Then  $u \wedge x = x$  and  $u \wedge y = y$



Now  $u \wedge (x \vee y) = (u \wedge x) \vee (u \wedge y) = x \vee y$

Therefore  $x \vee y \leq u$ , Hence  $\sup \{x, y\} = x \vee y$ .

**Theorem 3.9:** Let  $x, y \in A$  and  $x \wedge y = y \wedge x$  then  $\inf \{x, y\} = x \wedge y$

**Proof:** Let  $x, y \in A$  and  $x \wedge y = y \wedge x$ . Clearly,  $x \wedge y \leq x$

$$y \wedge (x \wedge y) = y \wedge (y \wedge x) = y \wedge x = x \wedge y$$

Hence  $x \wedge y \leq y$ . Therefore  $x \wedge y$  is a lower bound of  $\{x, y\}$

Suppose  $l$  is a lower bound of  $\{x, y\}$

$$x \wedge l = l = y \wedge l. \text{ Now } (x \wedge y) \wedge l = x \wedge (y \wedge l) = x \wedge l = l$$

Therefore  $l \leq x \wedge y$  and hence  $x \wedge y$  is the greatest lower bound of  $\{x, y\}$

**Corollary 3.10:** Let  $A$  be with  $T$ ,  $x \in B(A)$ ,  $y \in A$  such that  $x \wedge y = y \wedge x$  then

$x \vee y$  is the lub of  $\{x, y\}$  and  $x \wedge y$  is the glb of  $\{x, y\}$

Thus we have proved the following

**Theorem 3.11:** Let  $A$  be a C-algebra with  $T$ . Then the centre  $B(A)$  of  $A$  (def 1.12) is a Boolean algebra under the partial ordering ' $\leq$ ' induced from  $A$ .

Now we give the following equivalent conditions for " $x \leq y$ "

**Theorem: 3.12:** In a C-algebra

$$(i) \quad x \leq y \Leftrightarrow y \wedge (y' \vee x) = x \Leftrightarrow (y' \vee x) \wedge y = x$$

$$(ii) \quad x \leq y \Rightarrow y \vee x = y$$

**Proof:** (i) is clear from 2.3 and 2.4

$$(ii) \quad x \leq y \Rightarrow y \wedge x = x \Rightarrow y \vee (y \wedge x) = y \vee x \Rightarrow y = y \vee x$$

The converse of 3.12 (ii) is not true, for example, in 'C',  $U \vee F = U$  but  $F \not\leq U$

**Theorem 3.13.** In the poset  $(A, \leq)$

$$1. \quad x \vee y \leq x \vee x'$$

$$2. \quad \text{if } x \leq y \text{ then, for any } z \in A, (i) \ z \wedge x \leq z \wedge y \text{ (ii) } z \vee x \leq z \vee y$$

$$3. \quad x \wedge y \leq x$$

4. if A is a C-algebra with T and  $x \in B(A)$  then  $y \leq x \vee y$

**Proof:** 1.  $(x \vee x') \wedge (x \vee y) = (x \wedge (x \vee y)) \vee (x' \wedge (x \vee y))$   
 $= x \vee (x' \wedge y) = x \vee y$  therefore  $x \vee y \leq x \vee x'$

2. Suppose  $x \leq y$ , and  $z \in A$ .

(i)  $(z \wedge y) \wedge (z \wedge x) = (z \wedge y \wedge z) \wedge x = z \wedge y \wedge x$  (By 1.6)  
 $= z \wedge x$  (since  $y \wedge x = x$ )

Therefore  $z \wedge x \leq z \wedge y$

(ii) is a direct verification

3.  $x \wedge (x \wedge y) = (x \wedge x) \wedge y = x \wedge y$ . Therefore  $x \wedge y \leq x$ .

4. Suppose A is C-algebra with T and  $x \in B(A)$ . Then

$$(x \vee y) \wedge y = (x \wedge y) \vee (x' \wedge y \wedge y) = (x \wedge y) \vee (x' \wedge y)$$

$$= (x \vee x') \wedge y = T \wedge y = y$$

Therefore  $y \leq x \vee y$ .

### Note 3.14

(i) In the above,  $x \vee y \leq y \vee y'$  not hold; for example  $T \vee U \not\leq U \vee U'$  in C.

(ii) In the above,  $x \leq x \vee y$  not hold; for example  $F \not\leq F \vee U$  in C.

Now we prove modular type results in the following

**Lemma: 3.15:** In the poset  $(A, \leq)$

$$x \leq y \Rightarrow x \vee (y \wedge z) = y \wedge (x \vee z)$$

**Proof:** Suppose  $x \leq y \Rightarrow y \wedge x = x$

$$y \wedge (x \vee z) = (y \wedge x) \vee (y \wedge z) = x \vee (y \wedge z).$$

**Theorem 3.16:** For any  $x, y, z, \in A$

$$x \leq y, x \vee z = y \vee z \text{ and } x \wedge z = y \wedge z \text{ then } x = y$$

**Proof:**  $x = x \vee (x \wedge z) = x \vee (y \wedge z)$   
 $= y \wedge (x \vee z)$  (By above Lemma 3.15)

$$= y \wedge (y \vee z) = y$$

If  $x \in B(A)$  then we have proved (3.12) that  $y \leq x \vee y$  and if  $x \in B(A)$ ,  $x \wedge y = y \wedge x$  then  $x \vee y$  is  $\sup \{x, y\}$  (3.8). In general  $x \vee y$  need not be an upper bound of  $\{x, y\}$  in the poset  $(A, \leq)$ . If  $x \vee y$  is an upper bound of  $\{x, y\}$  in  $(A, \leq)$  then  $A$  becomes a Boolean algebra. Also it can be observed that, for  $x, y \in A$   $x \wedge y \leq x$  is true But  $x \wedge y \leq y$  need not hold in  $A$ . If  $x \wedge y \leq y$  for all  $x, y \in A$  then  $A$  becomes a Boolean algebra. We conclude this paper with the following.

**Theorem: 3.17:** The following are equivalent in a C-algebra  $A$

- (1)  $\langle A, \vee, \wedge, ' \rangle$  is a Boolean algebra
- (2)  $y \leq x \vee y$  in  $(A, \leq)$  for all  $x, y \in A$
- (3)  $x \leq x \vee y$  in  $(A, \leq)$  for all  $x, y \in A$
- (4)  $x \vee y$  is an upper bound of  $\{x, y\}$  in  $(A, \leq)$  for all  $x, y \in A$
- (5)  $\text{Sup } \{x, y\} = x \vee y$  in  $(A, \leq)$  for all  $x, y \in A$
- (6)  $x \wedge y \leq y$  in  $(A, \leq)$  for all  $x, y \in A$
- (7)  $x \wedge y$  is a lower bound of  $x, y$  for all  $x, y \in A$
- (8)  $\text{inf } \{x, y\} = x \wedge y$  in  $(A, \leq)$  for all  $x, y \in A$
- (9) The poset  $(A, \leq)$  is directed below
- (10)  $x \vee x'$  is the greatest element in the poset  $(A, \leq)$  for all  $x \in A$
- (11) For any  $x, y \in A$ , if  $y' \leq x'$  then  $x \leq y$

**Proof:** (1)  $\Rightarrow$  (2) is clear

(2)  $\Rightarrow$  (3): Assume (2), Therefore  $(x \vee y) \wedge y = y$  for all  $x, y \in A$

$$(x \vee y) \wedge (y \vee x) = (x \vee (y \vee x)) \wedge (y \vee x) \text{ (By 1.6 b)} = y \vee x$$

$$\text{Similarly } (y \vee x) \wedge (x \vee y) = x \vee y$$

Therefore  $x \vee y = y \vee x$ , for all  $x, y \in A$

Hence by (2)  $x \leq y \vee x = x \vee y$  for all  $x, y \in A$

(3)  $\Rightarrow$  (4): Assume, (3) we have  $(x \vee y) \wedge x = x$ , for all  $x, y \in A$

$\Rightarrow x \vee (y \wedge x) = x$  for all  $x, y \in A$ . By (2.8)

$$\begin{aligned} (x \wedge y) \vee (y \wedge x) &= (x \wedge y) \vee (y \wedge (x \wedge y)) \text{ (by 1.6b)} \\ &= x \wedge y \end{aligned}$$

Similarly,  $(y \wedge x) \vee (x \wedge y) = y \wedge x$

Therefore  $x \wedge y = y \wedge x$  for all  $x, y \in A$

Therefore  $x \vee y = y \vee x$  for all  $x, y \in A$

Now  $(x \vee y) \wedge y = (y \vee x) \wedge y = y$ , that is  $y \leq x \vee y$ .

$x \vee y$  is upper bound of  $\{x, y\}$  for all  $x, y$

(4)  $\Rightarrow$  (5): Assume (4). Suppose  $u$  is upper bound of  $\{x, y\}$ ,

$$u \wedge x = x, u \wedge y = y \text{ now } u \wedge (x \vee y) = (u \wedge x) \vee (u \wedge y) = x \vee y$$

Therefore,  $\sup \{x, y\} = x \vee y$  for all  $x, y \in A$

(5)  $\Rightarrow$  (6): Assume (5). Then,  $x \vee y = y \vee x$  for all  $x, y \in A$

so that  $x \wedge y = y \wedge x$  for all  $x, y \in A$

Now  $y \wedge (x \wedge y) = y \wedge (y \wedge x) = y \wedge x = x \wedge y$  Therefore  $x \wedge y \leq y$

(6)  $\Rightarrow$  (7): Assume (6). we know that  $x \wedge y \leq x$ .

Therefore  $x \wedge y$  is lower bound of  $\{x, y\}$ , for all  $x, y \in A$

(7)  $\Rightarrow$  (8): Assume (7). suppose  $t$  is lower bound of  $\{x, y\}$

$$x \wedge t = t = y \wedge t, \text{ Now } (x \wedge y) \wedge t = x \wedge (y \wedge t) = x \wedge t = t$$

Therefore  $x \wedge y$  is greatest lower bound of  $\{x, y\}$

(8)  $\Rightarrow$  (1): Assume (8). Therefore  $x \wedge y = y \wedge x$  for all  $x, y \in A$

Hence, by theorem 1.11,  $A$  is Boolean algebra

Therefore (1) to (8) are equivalent

Clearly (9), (10), (11) are equivalent to (1) by theorem 1.11.

Hence the theorem.

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